QoS enhancements for UHR-follow-up

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Introduction

- In 11be, some members highlighted the scalability problem with SCS since multiple flows are mapped to same TID (e.g., 11-20-1686) leading to possible HOL blocking problem at layer-2.
- To recap, the HOL blocking problem at layer-2 arises because an older packet waiting retransmission and mapped to a given TID would prevent newer packets that have been correctly received at layer-2 to be forwarded up.
 - This is not an issue where the newer packet belongs to same traffic flow AND at the application/transport layer the processing of the newer packet has dependence on that of the older packet (i.e., expect in-order delivery).
 - However, for cases when that's not true (see next slide), HOL blocking may affect userexperience by preventing chances of faster processing.
- Previously in 11-23-0697-uhr we proposed a solution to HOL blocking problem .
- In this document we provide further details on motivation and <u>evaluations.</u>

Use-cases overview

- There are situations where • Server-2 different traffic streams could (IP-2, port-2) be mapped to same TID Server-3 between two STAs such as Server-1 (IP-3, port-2 multiple voice/ video sessions, (IP-1, port-1) time-sensitive flows, file transfer instances etc. PPDU (TID-x) In some cases, there maybe ٠ additional information (e.g., **IP tuple, Port** PPDU (TID-y) #, metadata) available at
 - layer-2 to identify the different streams.

Use-cases overview (contd.)

- In other cases layer-2 may not have enough information to demultiplex streams because its encapsulated at higher layer (e.g., QUIC connection consisting of multiple streams).
- In addition, there are situations where partial delivery of packets belonging to same flow is also useful. For example,
 - when application layer FEC [1] is applied in which case the receiver may make use of whatever packets are received.
 - An older packet that arrived after a deadline and the newer frames are not dependent on the old packet.
 - Potential packet loss concealment techniques in webrtc that uses AI to fill in missing packets [3]

Demultiplexing streams at layer-2

- For the case when the transmitter STA (e.g., an AP) has enough information to demultiplex streams and each of those streams require in-order ("IO") delivery (e.g., TCP), the HOL blocking can be resolved if we can put them on to different TIDs.
- Clearly, this is a problem if the number of such streams exceed the number of TIDs that can today be mapped to an AC.
- Proposal: the UHR spec can allow more than 2 TIDs to be mapped to an AC when needed.

Mapping multiple TIDs to an AC

- Today typically TID = UP and UPs are restricted from 0-7. Also, there is explicit mapping from those 8 UPs to 4 ACs.
- Option 1: keep current UP to AC mapping but define mapping from TIDs 8-14 to an AC.
 - Only subset of TIDs 8-14 may be mapped and only when there is a flow that needs it (e.g., during an active SCS stream).
- Option 2: allow a temporarily unused TID between 0-7 to be mapped to a different AC than that's derived from default mapping.
 - Could be dynamic as well (i.e., only during an active SCS stream).

Allowing out-of-order delivery (recap from 11-23-0697)

- For the cases when demultiplexing streams at layer-2 is not feasible but possible at upper layer (e.g., QUIC, RTP etc.) and when the upper layer can handle some packet reordering, we propose to allow the MAC-SAP to deliver out-of-order ("OOO") packets optionally for a subset of TIDs.
 - Maybe limit to 1 or 2 TIDs. Also, may use only when a flow that can benefit from OOO delivery is established (i.e., following a corresponding SCS negotiation).
- This is similar to out-of-order delivery by PDCP feature in 3GPP.
- Frame replay detection:
 - Today the in-order delivery sequence allows for easy frame replay detection (i.e., just check if the PN number is greater than last highest PN number).
 - For out-of-order delivery, the frame replay detection needs to change (e.g., by maintaining a sliding scoreboard of PNs).
 - Need a way to separate the PN space used by frames requiring OOO delivery vs ones that require IO.

Evaluations (Scenario 1)

• Network:

- 1 DUT (i.e., wifi-8) AP with one DUT STA and DL flow. Single link, BW = 40 MHz, MCS-5, NSS-2.
- N OBSS APs (where N = 0, 1, ...5).
- TXOP at DUT and OBSS AP = 6 ms.

• Traffic and KPI:

- At DUT, model very simplistic HTTP response traffic
 - generate N objects of size M bytes one time where set of $(N, M) = \{(5, 80K), (10,40K), (20,20K)\}$.
 - Interleaved packets modeling network jitter/delay.
 - Measure latency for each IP packet of an object as well overall latency for each object.
 - For example, this could model progressively loading different parts of a webpage.
- OBSS traffic is full buffer.
- Link level PER: (1%, 10%).

Results (scenario 1)



Results (scenario 1 contd.)



Results (scenario 1 contd.)



Similarly, overall latency for the whole object also improves significantly (e.g., from 77ms to 55ms when number of OBSS STAs is 5, number of objects is 20 and PER is 10%).

Evaluations (Scenario 2)

• Network:

- 1 DUT AP with one DUT STA and DL flow. Single link, BW = 80 MHz, MCS-7 or 6, NSS-2.
- N OBSS APs (where N = 3, 5).
- TXOP at DUT AP = 3 ms.

• Traffic and KPI:

- At DUT, 4 flows modelling 4 point cloud objects [2]
 - Flow-1 to 4 has data rates of 53.5, 21.95, 13.63 and 10.0 Mbps respectively with corresponding video frames generated every 16ms and split into 1400B IP packets.
 - Interleaved packets modeling network jitter/delay.
 - Measure latency for each IP packet of a frame of an object as well overall latency for each frame.
- Each OBSS AP occupy X % of airtime (where X = 10, 20).
- Link level PER: (1%, 10%).

Results (scenario 2 with MCS-6)



Clear gains until the network below saturation, then the DUT has difficulty delivering all traffic.

Results scenario 2 average latency per IP packet per frame



- Significant reduction in latency (e.g., about 20-30% reduction in latency with number of OBSS =3, PER = 10%).
- Low data rate flows (e.g., flow-4) benefits more because they are more likely to be blocked by a packet from a high data rate flow.

Results scenario 2 average latency per frame



• Also, significant reduction in overall frame latency esp. for low data rate flows (~ 10%).

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Possible high-level design to enable OOO delivery

- Limit to changes only at driver/software level. No change in low level MAC design including BA score-boarding.
- No mixing of IO and OOO delivery in same TID at any time.
- Separate key generated during 4-way handshake used to encrypt TIDs with OOO traffic => PN space for OOO frames are different from that of IO frames.
- No change in reception for TIDs that contain flows with IO traffic.
- The PN window should be equal to size of the maximum number of MPDUs for that TID that can be aggregated in a single PPDU i.e., BA scoreboard size for that TID.
 - Otherwise, we will have scenarios where a retransmitted packet wont be forwarded up for failing the replay detect.
 - Also, transmitter wont send frames beyond the scoreboard length.

Summary

- Proposed couple of solutions to resolve the HOL blocking issue at MAC layer.
- Also, presented some simulation results showing value of resolving this.

References

1. 3GPP TR 26.926 V18.1.0 (2023-12)

- 2. 8I Voxelized Full Bodies—A Voxelized Point Cloud Dataset
- 3. https://blog.research.google/2020/04/improving-audio-quality-in-duo-with.html

Results (scenario 1)



Results (scenario 2 with MCS-6)



Clear gains until the network below saturation, then the DUT has difficulty delivering all traffic.

Results (scenario 2 contd. With MCS-7)



Clear gains until the network is below saturation then the DUT has difficulty delivering all traffic.

Submission