IEEE P802.11  
Wireless LANs

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| MLD architecture part 2 | | | | |
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Abstract

This submission builds upon the MLD architecture in TGbe D1.31 (as incorporated from 11-21/0577, by adding support for group addressed transmissions from AP MLD, and legacy operation.

R0 – Initial discussion document.

R1 – Added proposed resolution text for the relevant CIDs.

R2 – After discussion of Affiliated STA definition, captured alternate concept that an affiliated STA/AP is defined to provide the communication path between MLDs, without further details.

R3 – Agreed the non-AP STA is either MLD or “legacy” and never a combination (as opposed to the AP MLD, which is a combination). Wording needs to be updated to align with this agreement. Clarified power save state tracking versus power save buffering, w.r.t. upper/lower MAC sublayer functions.

R4 – Updated based on comments from Aug 30 teleconference. Updates are labelled with “8/30”.

R5 – Updates during Sept 9 teleconference (marked with “9/9”).

R6- Updates follow Sept 9 teleconference, based on discussion on that call. Marked with “9/13”.

R7 – Updates per 11-21/0209 review.

R8 – Updates from Sept 14 telecon. Marked with 9/14.

R9 – Updates from Sept 15 telecon. Marked with 9/15.

R10 – Significant updates and simplification, from off-line discussion.

R11 – Removed that affiliated STAs are co-located with the MLD entity. Fixed power save details in 5.1.5.1a bullet list (functions in different entities). Editorial change to vertically align affiliated APs in Figure 4-30c (to match Figure 5-2b).

R12 – Fixed file header and date.

R13 – Corrected the AP MLD requiring “more than one” affiliated AP, to align with 11-21/2009 that it can sometimes operate with only one affiliated AP. Removed most uses of “non-MLD”, especially applied to a non-AP device - just say “non-AP STA” to mean a non-MLD/legacy client. Also reworded the concept of a non-AP MLD switching “modes” to be a legacy (non-MLO) STA, to remove the concept of “mode”.

**Introduction**

This document extends the MLD architecture concepts in the TGbe draft, to add architectural support for group addressed frames and legacy operation (of affiliated STAs), along with some (mostly) editorial suggestions.

**CC36 CIDs:**

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| 5171 | Guogang Huang | 3.2 | 41.16 | For the AP MLD, If there are legacy STAs which associate with affiliated APs, then each affiliated AP will have a MAC SAP to LLC, which is identify by the MAC address of the corresponding affiliated AP. | Please add a note below the MLD definition, e.g.  Note. For an AP MLD, If there are legacy STAs which associate with each affiliated AP, then each affiliated AP will also have a MAC SAP to LLC, which is identify by the MAC address of the corresponding affiliated AP. | **Revised.**  **The MAC SAP at an AP serves the DSAF, not LLC. However, the intent of the clarification is agreed.**  **Make the changes shown in 11-21/1111, which adds the architectural structure for affiliated APs to the AP MLD concepts.** |
| 5172 | Guogang Huang | 4.9 | 49.44 | Add a subclause 4.9.5 to describe the reference model for MLD and explain the legacy support of the AP MLD | As in comment. | **Revised. Make the changes shown in 11-21/1111, which adds the architectural structure for affiliated APs to the AP MLD concepts.** |
| 5173 | Guogang Huang | 7.1 | 0.00 | Update the figure 7-1 DS architecture, and clarify the number of DS SAPs for an AP MLD especially when there are legacy STAs associated with each affiliated AP | As in comment. | **Revised. Make the changes shown in 11-21/1111, which adds Figure 7-2 to show the DS SAPs, as requested.** |
| 6187 | Michael Montemurro | 4 | 45.01 | In MLO, affiliated APs are able to provide BSS connectivtiy to legacy STAs but there is no description on how this works. Proide a description of how an affiiated AP can service legacy STAs while also operating with an MLD to support MLO. | The commenter is willling to collaborate on a contribution which would add a description to address this comment. | **Revised. Make the changes shown in 11-21/1111, which adds the architectural structure for affiliated APs to the AP MLD concepts.** |
| 7349 | Stephen McCann | 3.2 | 41.12 | A definition of "affiliated" would be useful in clause 3.2 | Add the following definition "Affiliated: A STA and an MLD that are co-located or connected through an existing security relationship." | **Revised. Make the changes shown in 11-21/1111, which adds a definition for affiliated** |

**Discussion:**

1. **Group addressed MSDU handling and legacy operation:**

**Group addressed MSDU handling:**

11-21/0577 (and therefore the current draft text in subclause 4.9 and clause 5) touched on group addressed frames, but only superficially. It mentioned that there are separate GTK/IGTK/BIGTK maintained per link, and that group addressed frame encryption and decryption are done with the GTK for a given link. Further clarification and details are needed, however.

First, it needs to be noted that with separate group keys per link, there needs to be separate key management per link. This implies a separate Authenticator/Supplicant “context” per link paired with the link’s specific RSNA key management for group keys. Thus, our architectural model for an MLD (Figure 4-30a) needs to be extended to add these RSN facilities per link, for the group addressed frames and group key handling.

Note that these RSN facilities are the same ones as used by the affiliated APs for legacy operations (with associated non-MLD peers). See immediately next section…

**Legacy AP operations:**

Legacy operation of the affiliated APs is a key feature, which raises some architectural details. In particular, as upper MAC operations for MLD peers is modified from single link (legacy) operations, it is simplest to model the AP MLD and its affiliated legacy APs as having separate upper MAC facilities. Per Figure 4-30a and Figure 4-30b in the draft, the AP MLD’s upper MAC facilities are performed by a single MLD upper MAC sublayer which operates using multiple links and multiple MLD lower MAC sublayer entities. However, each affiliated AP provides legacy upper MAC operations to associated legacy non-AP STAs, through an independent MLD upper MAC sublayer. A new figure (Figure 4-30c, below) is added to show this architecture of an AP MLD and its (legacy supporting) affiliated APs.

Since the MLD lower MAC sublayer is shared between the MLD and legacy upper MAC sublayers, the operation of this sharing needs to be specified where it affects externally visible behavior. To support group addressed security contexts (one set of GTK/IGTK/BIGTK per link), and legacy operation (PTK per peer STA) an AP MLD maintains an Authenticator per link (with legacy association PTKSAs, as well as GTK/IGTK/BIGTK for that link) and an Authenticator for the AP MLD (with PKTSAs per MLD association). These Authenticators cooperate as needed, on key updates, etc.

**Non-AP MLD and legacy non-AP STA:**

On the non-AP MLD, there is no need for this complexity. A non-AP MLD is simply a MLD upper MAC sublayer, and multiple MLD lower MAC sublayers (one for each link). Within the non-AP MLD, the combination of a single link lower MAC functionality plus the shared upper MAC results in a complete affiliated STA. A new figure (Figure 4-30d, below) is added to show the architecture of a non-AP MLD and its affiliated STAs. A non-AP device that switches from MLO to legacy operation (when roaming to a legacy AP, for example), retains the upper MAC sublayer and only one lower MAC sublayer, thus switching to operation over a single link, but making the transition invisible to the upper layers. Further, only one Supplicant is modelled in the reference model for the non-AP MLD. However, this Supplicant is enhanced over a legacy non-AP STA, in that it can manage the multiple sets of GTK/IGTK/BIGTK, one set per each link.

1. **Minor updates to Figure 5-2a, new Figure 5-2b**

Figure 5-2a is modified editorially. The light grey boxes in the draft, showing the upper and lower MAC functionalities will not copy well and the layers of boxes get confusing, so suggest replacing with braces instead. Also, the separation of TX and RX functions (left side and right side of the stack) in the upper MAC is carried over into the lower MAC, for consistency. Lastly, the space between the left (transmitting) and right (receiving) sides of the stack is removed, as this space is creating what looks like an unlabelled “box” in the architecture.

An additional figure and text are added following Figure 5-2a, to introduce the MPDU distribution function at the top of the MLD lower MAC sublayer to support legacy operation (distributing to the MLD or an affiliated AP, based on a mapping from the TA) and the transmission of group addressed frames.

1. **Function(s) of MLD lower MAC sublayer:**

As the MLD lower MAC sublayer comprises only some link-specific, low-level MAC functions (as shown in Figure 5-2a), this component cannot perform many complex MAC functions. Instead, for an AP MLD, it is actually the non-MLD upper MAC sublayer of an affiliated AP that performs these functions.

The discussion in 4.9.5 of where functions are performed needs to be updated to clarify this. Examples include:

* Maintenance of Link-specific GTK/IGTK/BIGTK (between an AP affiliated with the AP MLD and a STA affiliated with the non-AP MLD)
* Link-specific encryption/decryption/integrity protection and PN assignment using GTK/IGTK/BIGTK (between an AP affiliated with the AP MLD and a STA affiliated with the non-AP MLD)

And, some functions described in 4.9.5 (of 11-21/0577) could be clarified that it is only the *tracking* of this information that is (or could be) performed in the MLD lower MAC sublayer, including:

* Power save state and mode

while power save queuing and timing are managed by the MLD upper MAC sublayer.

1. **“Affiliated”**

Multiple places in the TGbe draft refer to affiliated STAs/APs (of course). But, this term is never clearly defined. We do have the following definitions (TGbe draft) which imply the concept:

**access point (AP) multi-link device (MLD):** An MLD, where each station (STA) affiliated with the MLD is an AP.

**multi-link device (MLD):** A device that is a logical entity and has more than one affiliated station (STA) and has a single medium access control (MAC) service access point (SAP) to logical link control (LLC), which includes one MAC data service.

**non-access point (non-AP) multi-link device (MLD):** An MLD, where each station (STA) affiliated with the MLD is a non-AP STA).

Concepts related to “affiliated”:

* When an MLD association is done, the affiliated APs/STAs provide the communication paths (links) between the MLDs;
* An affiliated AP provides the pre-association services for an AP MLD;
* The “lower” services of the MAC and link-specific behaviors are described as operations of the affiliated STA/AP, throughout the changes to clauses 10 and 11, and the multi-link operation in 35.3 (and in other places, as well.

So, an affiliated STA/AP seems to provide the lower MAC services for an MLD (and possibly some higher MAC services such as encryption in some scenarios).

However, there are many examples of phrasing such as, “A STA affiliated with an MLD shall/may …” So, it seems that these “affiliated” entities are themselves complete STAs (AP or non-AP).

Since significant parts of the MAC operation of an MLD are non-link-specific, and shared in a single MLD entity that uses/controls all the links, clearly the “upper” services of the MAC that operate at the MLD level are required to complete the STAs of an “affiliated” STA.

Thus, it seems the ‘complete’ stack is the logical view, including both the “lower” MAC/link-specific services for one link (and one PHY) and an “upper” MAC component that performs MLD level operations.

This leads to a definition of affiliated as something similar to:

**affiliated**: A STA (AP or non-AP) that provides link-specific, lower MAC services within an MLD. .

1. **Reorganize text from 11-21/0577**

Now that the concepts for MLD are complete (including the “legacy” operations), we can re-visit the organization of the introductory material in clause 4. There are generally three types of material: 1) an introduction to the concepts of multi-link operation between peers; 2) some description, still high-level/reference model, to expand on these architectural concepts at the level of the rest of clause 4, just clarifying concepts but not operational structure and details; 3) details that do get into the operation of the concepts and the detailed components of the stack.

Recommendation: Arrange the text in the new subclause 4.9.5 by introducing the multi-link concepts at a high-level first (the text shaded in red below), and then the reference model that can support this behavior (the text shaded in blue below). Move text shaded in green from clause 4 into clause 5, as this introduces the details of stack components and helps explain Figure 5-2a, and is more relevant where those details are discussed.

**Proposed Changes:**

***Changes proposed by this document are shown with underscore*** ***additions and ~~strikethrough~~ deletions, with TGbe D1.31 as baseline.***

***Track changes are used for discussion purposes only, as revisions of this document are reviewed and edited.***

***TGbe editor: Please add a definition to subclause 3.2: :***

**affiliated**: A STA (AP or non-AP) that provides link-specific, lower MAC services within an MLD.

***TGbe editor: Please add the unshaded text indicated below Note that text shaded in red or shaded in blue below is existing text from TGbe D3.1, rearranged in order, and with some edits as indicated. The text shaded in green is moved from clause 4 to clause 5, and with some edits as indicated.***

**4.9.5 Reference model for multi-link operation (MLO)**

MLO allows operation over multiple links. An MLD manages such communication over multiple links. Communication across different frequency bands/channels can occur simultaneously or not depending on the capabilities of both the AP MLD and the non-AP MLD (see 35.3.13.2 (Simultaneous transmit and receive (STR) operation) and 35.3.13.3 (Nonsimultaneous transmit and receive (NSTR) operation)).

The MLO procedures (see 35.3 (Multi-link operation)) allow a pair of MLDs to discover, synchronize, (de)authenticate, (re)associate, disassociate, and manage resources with each other on any common bands or channels that are supported by both MLDs.

As described in 35.3.1 (General), each AP MLD has a single MAC-SAP and each non-AP MLD has a single MAC-SAP. Each AP affiliated with an AP MLD has a ~~different~~ MAC address different from any other AP affiliated with the AP MLD, ~~within the MLD~~ and each STA affiliated with a non-AP MLD has a ~~different~~ MAC address different from any other STA affiliated with the non-AP MLD~~within the MLD~~.

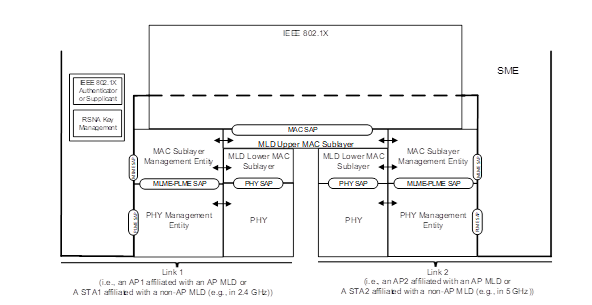
An example of an AP MLD with two links (Link 1 and Link 2) is shown in Figure 4-29b (Example MLD and the affiliated STA communication system). The figure shows a~~A~~n AP MLD with MLD MAC address *M* and the MLD lower MAC sublayers of two affiliated APs (AP1 with MAC address *w* and AP2 with MAC address *x*). The AP MLD is associated with a non-AP MLD with MLD MAC address *P* and the MLD lower MAC sublayers of two affiliated STAs (STA1 with MAC address *y* and STA2 with MAC address z*)* are shown*.* Link 1 is established between AP1 and STA1 and link 2 is established between AP2 and STA2. In general, the MAC address of an MLD and the MAC addresses of the STAs affiliated with the MLD are all different (e.g., *M*, *P*, *w*, *x*, *y*, and *z* have different values).



**Figure 4-30b – Example MLD and the affiliated STA communication system**

The reference model of a multi-link device (MLD) (see 35.3 (Multi-link operation)) is shown in Figure 4-29a (Reference model for an MLD).

NOTE—For simplicity, Figure 4-29a (Reference model for an MLD) depicts the reference model when there are two links, while in general, an MLD can support more than two links.



**Figure 4-30a – Reference model for an MLD for two links**

NOTE—The SME boundary top is left open in Figure 4-29a (Reference model for an MLD) to indicate that the SME can contain other functions that are not defined by this standard.

An MLD supports multiple MAC sublayers, coordinated by an SME.

The SME maintains the authentication and association states. The Authenticator and the MAC-SAP of the AP MLD are identified by the same AP MLD MAC address. The Supplicant and the MAC-SAP of the non-AP MLD are identified by the same non-AP MLD MAC address.

The SME is responsible for coordinating each of the MLMEs of all affiliated STAs, and to maintain an ~~single~~ RSNA key management entity~~, as well as a single~~ and IEEE 802.1X Authenticator or Supplicant in each upper MAC sublayer component, for MLO.

The MAC Sublayer is further divided into an MLD ~~U~~upper MAC sublayer and an MLD ~~L~~lower MAC sublayer. The MLD ~~U~~upper MAC sublayer ~~(MLD)~~ performs functionalities that are common across all links, and the MLD ~~L~~lower MAC sublayer (shared with an AP or STA affiliated with the MLD) performs functionalities that are local to each link. Some of the functionalities require joint processing of both the MLD ~~U~~upper MAC sublayer and MLD ~~L~~lower MAC sublayer~~s~~.

An AP MLD always operates in cooperation with one or more affiliated APs, one for each physical link. The MLD lower MAC sublayer components implement link-specific functions which operate independently of the lower MAC in other affiliated APs, and are shared between each affiliated AP and the AP MLD operations. Some behaviors of MLO require the use one or more affiliated APs’ upper MAC components. In particular, the affiliated AP MLD upper MAC sublayer components support group addressed traffic, and traffic to or from any non-AP STAs (which are not capable of MLO). The high-level structure of an AP MLD along with its affiliated APs is shown in Figure 4-30c.



**Figure 4-30c – High-level architecture for AP MLD with affiliated APs**

The non-AP MLD reference model includes the MLD upper MAC sublayer and MLD lower MAC sublayers (one for each link). The single upper MAC within a non-AP MLD can operate at any given time in either MLO over multiple lower MAC and PHY pairs for association to an AP MLD, or as a (non-MLO) non-AP STA using only one set of lower MAC and PHY for association to an AP that is not an affiliated AP. A single Supplicant on the non-AP MLD manages the PKTSA, and multiple group key security associations (one set per link). The reference architecture is shown in Figure 4-29d.



**Figure 4-30d – High-level architecture for non-AP MLD with affiliated STAs**

***TGbe editor: Please modify subclause 5.1.5.1a as follows. Note that the text shaded in green is moved from clause 4 to clause 5, and with some edits as indicated***

**5.1.5 MAC data service architecture**

**5.1.5.1a Multi-link Operation**

For Multi-link Operation (MLO), one or more links are used for communication between the AP MLD and non-AP MLD after MLD (re)setup as described in 35.3.5 (Multi-link (re)setup)). The MAC data plane architecture of an MLD with *n* links (i.e., processes that involve transport of all or part of an MSDU) is shown in Figure 5-2a (MAC data plane architecture (MLO)).

~~~~



**Figure 5-2a - MAC data plane architecture (MLO) for unicast data frames**

During transmission, an MSDU from the MAC SAP goes through the processes shown in the left-hand side of Figure 5-2a (MAC data plane architecture (MLO)), then through the TID-to-link mapping process (see 35.3.6.1 (TID-to-link mapping)) that forwards the MPDUs down to one of the MLD ~~L~~lower MAC sublayers and then to the corresponding PHY SAP.

Note – TID-to-link mapping negotiation between peer MLDs is an optional feature.

During reception, MPDUs originating from different PHY SAPs first go through an MLD ~~L~~lower MAC sublayer, followed by a merging process, and then go through the rest of the processes in the right-hand side of Figure 5-2a (MAC data plane architecture (MLO)). Then, one or more MSDUs are delivered to the MAC SAP or, via the DSAF to the DS. The IEEE 802.1X Controlled/Uncontrolled Ports discard any received MSDUs if the Controlled Port is not enabled and if the MSDU does not represent an IEEE 802.1X frame.

NOTE—Many of the processes shown in Figure 5-2a (MAC data plane architecture (MLO)) also apply to MLD-level MMPDU flows for the MAC control plane architecture, and the processes shown at the MLD ~~L~~lower MAC sublayer also apply to Control and Extension frames.

When MLO is being used, the same security association (PTKSA) is used to encrypt the unicast MPDUs and MMPDUs prior to transmission on all the links. The same security association (PTKSA) is used to decrypt the unicast MPDUs and MMPDUs received on all the links.

For an AP MLD to support group address transmissions and also non-MLO peer STA associations, Figure 5-2a is combined with *n* affiliated APs, within a structure as shown in Figure 4-29c. The MLD upper MAC sublayer components of the affiliated APs are the same as those for the AP MLD, but handle group addressed security associations (GTK, IGTK and BIGTK), and handle traffic to and from associated non-AP STAs (not operating in MLO) with single-link security associations for peerwise keys (PTKs). The overall structure is as shown in Figure 5-2b.



**Figure 5-2b -** **MAC data plane architecture for MLD AP and affiliated APs**

An additional function is added for data MPDU reception to distribute the MPDUs to the appropriate MLD upper MAC sublayer based on the type of association with the peer, which is tracked per TA. MPDUs received from an MLD non-AP STA peer are delivered to the AP MLD upper MAC, and other MPDUs are delivered to the affiliated AP upper MAC for that link.

Group addressed MSDUs at the DS are not transmitted directly by affiliated APs. Instead, the MLD AP processes group addressed MSDUs to the point of assigning a sequence number. The MLD AP and affiliated APs then coordinate to power save buffer (if appropriate), assign packet numbers and encrypt the resulting MPDU in the individual affiliated APs’ stacks. The GTK of ~~a link~~ an affiliated AP is used to encrypt the group addressed frames MPDUs and MMPDUs prior to transmission on the link managed by that affiliated AP. The GTK of ~~a link~~ the corresponding affiliated STA is used to decrypt the group addressed frames MPDUs and MMPDUs received on ~~the~~ a link. Group-addressed MMPDUs generated within the AP MLD upper MAC sublayer shall be transferred to the appropriate affiliated APs for transmission.

NOTE--An implementation must confirm that an MSDU that would otherwise be transmitted to peer MLD STAs is still transmitted, even if group address filtering of multicast MSDUs is being performed such that the MSDU might not be transmitted by the affiliated APs.

The MLD ~~U~~upper MAC sublayer functions include:

* Authentication, association and reassociation (between an AP MLD and a non-AP MLD)
* Security association (e.g., PMKSA, PTKSA) and distribution of GTK/IGTK/BIGTK
* SN/PN assignment for frames to be encrypted by PTK for unicast frames
* Power save buffering of individually addressed frames (only on AP MLD)
* Encryption/decryption using PTK for unicast frames
* Selection of the MLD ~~L~~lower MAC sublayer for transmission (TID-to-link mapping (see 35.3.6.1 (TID-to-link mapping)))
* Reordering of packets to ensure in-order delivery per each BA session
* Block Ack scoreboarding for individually addressed frames (in collaboration with the MLD ~~L~~lower MAC sublayer). Optionally, the MLD ~~U~~upper MAC sublayer delivers the BA record on one link to the MLD ~~L~~lower MAC sublayer of other links)
* MLD-level management information exchange/indication via the MLD ~~L~~lower MAC sublayer

The affiliated (non-MLO) upper MAC sublayer functions (only on AP) include:

* Non-MLO peer operations, above the MLD lower MAC sublayer
* Maintenance of Link-specific GTK/IGTK/BIGTK (between an AP affiliated with the AP MLD and a STA affiliated with the non-AP MLD)
* Link-specific encryption/decryption/integrity protection and PN assignment using GTK/IGTK/BIGTK (between an AP affiliated with the AP MLD and a STA affiliated with the non-AP MLD)
* Link-specific management info exchange/indication (e.g., Beacon)
* Power save buffering of group addressed frames

The MLD ~~L~~lower MAC sublayer functions include:

* ~~Maintenance of Link-specific GTK/IGTK/BIGTK (between an AP affiliated with the AP MLD and a STA affiliated with the non-AP MLD)~~
* ~~Link-specific encryption/decryption/integrity protection and PN assignment using GTK/IGTK/BIGTK (between an AP affiliated with the AP MLD and a STA affiliated with the non-AP MLD)~~
* ~~Link-specific management info exchange/indication (e.g., Beacon)~~
* Link-specific control info exchange/indication (e.g., RTS/CTS, Acks, NDP, etc.)
* Power save state and mode
* MAC address filtering for frame reception
* Block Ack scoreboarding for individually addressed frames (in collaboration with the MLD ~~U~~upper MAC sublayer). Optionally, the MLD ~~L~~lower MAC sublayer receives ~~from the~~ the BA record on the other links from the MLD ~~U~~upper MAC sublayer)

NOTE—The above functionality partitioning is meant for modelling the functionalities of each MAC Sublayer and is not meant for describing the MAC Sublayer for which the actual implementation of each function should reside.

NOTE – The Block Ack scoreboarding maintenance collaborated between the MLD ~~U~~upper MAC sublayer and MLD ~~L~~lower MAC sublayer is implementation dependent.

When MLO is being used, the “Block Ack Scoreboarding” block in the MLD ~~U~~upper MAC sublayer manages the BA status of the MPDUs (of this BA session) that are received on any setup link. The “Block Ack Scoreboarding” block in the MLD ~~L~~lower MAC sublayer manages the BA status of the MPDUs (of this BA session) that are received on this link. It may convey BA status of the MPDUs received on another link if it obtained such info from the other link via the MLD ~~U~~upper MAC sublayer.

***TGbe editor: Please modify this subclause as shown***

**5.1.5.11 AP MLD role**

In an AP MLD, the MAC data plane architecture as shown in Figure 5-2a (MAC data plane architecture (MLO) for unicast data frames) and Figure 5-2b (MAC data plane architecture for MLD AP and affiliated APs) includes Distribution System (DS) access in its role-specific behavior block, as shown in Figure 5-12 (Role-specific behavior block for an AP MLD). This block provides access to the DS for associated non-AP MLDs as described in 4.5.2.1 (Distribution).

NOTE—This behavior block indicates that there is no access through the controlled port to or from the local ~~upper~~higher-layers (e.g., the LLC sublayer) at an AP MLD. Any such access is logically achieved in the architecture via transition of the DS and Portal to an integrated LAN. In actual implementations, this is likely to be optimized, and Data frames appear to be delivered directly to one or more local LLC sublayer entities on the same physical device as the AP MLD. Such optimization is effectively distributing the functions of the DS and Portal, and it is the responsibility of the implementation to ensure the logical behavior of these entities is maintained.



**Figure 5-12 - Role-specific behavior block for an AP MLD**

***TGbe editor: Please modify subclause 7.1 as follows:***

**7.1 Introduction**

The DS SAP is the interface between the DS SAP service users and the DS SAP service provider. The DS SAP service users are the connected APs, mesh gates, the portal, and AP MLDs. The DS SAP service provider is the DS. Figure 7-1 (DS architecture(#2251)) shows the location of the DS in the IEEE 802.11 architecture. The DS SAP is indicated in this Figure by the lines connecting the DS to its service users. In Figure 7-1 (DS architecture(#2251)), the DS has four users, two APs, a mesh gate, a portal, and an AP MLD, so the DS is shown passing behind the MAC/PHYs of the STAs.



**Figure 7-1 – DS architecture**

The DS SAP interface specification describes the primitives required to get MAC service tuples in and out of the DS and

* update the DS’s mapping of STAs to APs or to mesh gates,
* update the DS’s mapping of non-AP MLDs to AP MLDs

Describing the DS itself or the functions thereof is out of scope of this standard.

The DS SAP actions are as follows:

1. Accept MSDUs (as part of MAC service tuples) from APs, mesh gates, the portal and AP MLDs.
2. Deliver MSDUs (as part of MAC service tuples) to APs, mesh gates, the portal, or the AP MLDs.
3. Accept STA-to-AP mapping updates from the APs.
4. Accept STA-to-mesh gate mapping updates from the mesh gates.
5. Accept non-AP-MLD-to-AP-MLD mapping updates from the AP MLDs.

NOTE—For MLDs, the source address or destination address parameters of the MAC service tuples (see 5.2.3.2 (Semantics of the service primitive)) are set to the MLD MAC address of the non-AP MLD, which is the identity of the non-AP MLD known by the DS.

When the DS delivers the MAC service tuples to an AP, the AP then determines when and how to deliver the MAC service tuples to the AP’s MAC (via the MAC SAP). When the DS delivers the MAC service tuples to a mesh gate, the mesh gate then determines when and how to deliver the MAC service tuples to the mesh gate’s MAC (via the MAC SAP). When the DS delivers the MAC service tuples to an AP MLD through DSAF, the AP MLD then determines when and how to deliver the MAC service tuples to the AP MLD’s MLD ~~U~~upper MAC sublayer (via the MAC SAP).

In the case of an AP MLD and its affiliated AP(s) connected to the DS, there are individual DS SAPs for each affiliated AP and one for the AP MLD, as shown in Figure 7-2. The affiliated AP(s) will each provide a mapping to their associated non-AP STAs, by their MAC addresses. The AP MLD will provide a mapping to its associated non-AP MLDs, by their MLD MAC addresses. Thus, the non-AP devices form distinct sets of MAC addresses, and the DS can deliver any service tuples with a one-to-one mapping of destination address to DS SAP.



**Figure 7-2 – Example DS access for an AP MLD with two affiliated APs**