TGaa OBSS Background

Date: 2009-07-10

Authors:

Name	Affiliations	Address	Phone	email
Graham Smith	DSP Group	2491 Sunrise Blvd, #100, Rancho Cordova, CA 95742	916 851 9191 X209	Graham.smith@dspg.com

Abstract

- TGaa has been working on a proposal for OBSS. This work has included much background investigation that is considered to be of interest to TGac.
 - What is the OBSS problem?
 - Sizing OBSS
 - Channel Selection
 - TGaa proposed solution "QLoad"
 - Sharing

doc.: IEEE 802.11-09/0762-00-00aa

What is the problem?

#	Network A	OBSS Network B	Effect	Result
1	Legacy	Legacy	Traffic simply competes	 Reduced bandwidth in each network No lost packets Not recommended for streaming
2	EDCA	Legacy	Higher priority traffic in Network A will drive down traffic in Network B	 AC_VO and AC_VI traffic dominates. Could be OK for streaming traffic but no admission policy Network A "wins"
3	EDCA	EDCA	Traffic competes on a priority basis. Networks compete on an 'equal' basis	 Reduced bandwidth in each network No real protection for streaming traffic in either network

Effects of OBSS - 1

#	Network A	OBSS Network B	Effect	Result
4	Admission Control	Legacy	Higher priority traffic in Network A will drive down traffic in Network B	 AC_VO and AC_VI traffic dominates. Could be OK for streaming traffic Network B bandwidth can be drastically reduced
5	Admission Control	EDCA	Traffic competes on a priority basis. Admission Control in Network cannot control traffic in Network B	• No protection for admitted traffic in Network A
6	Admission Control	Admission Control	Traffic competes on a priority basis. Admission Control in either Network cannot control traffic in other Network	• No protection for admitted traffic in either Network

Effects of OBSS - 2

These cases are cause for concern, Admission Control is intended to provide QoS 'protection', and it breaks down in OBSS!

7	НССА	Legacy	Scheduled TXOPs in Network A also apply CFP to Network B.	 Full protection for scheduled traffic in Network A Network B bandwidth reduced
8	НССА	EDCA	Scheduled TXOPs in Network A also apply CFP to Network B.	 Full protection for scheduled traffic in Network A Network B bandwidth reduced
9	HCCA	Admission Control	Scheduled TXOPs in Network A also apply CFP to Network B Admitted traffic Network B is lower priority than scheduled traffic in Network A	 Full protection for scheduled traffic in Network A Network B bandwidth reduced Both Networks using TSPECS
10	НССА	HCCA	Each HCCA AP will admit streams and allocate time to them BUT each AP and STA will obey the TXOP allocation of the other. No guarantee that each Network can allocate time when it needs to.	• Reduced protection for scheduled traffic in either network.

Effects of OBSS - 3

OBSS and **QoS**

- 1. For non-QoS (non-real time streaming) applications OBSS is simply a sharing or reduced bandwidth per network – Not a significant problem, if we assume DCF works!
- 2. OBSS is a significant problem ONLY when QoS is used AND when some 'guaranteed performance' is at stake

doc.: IEEE 802.11-09/0762-00-00aa

How big is the problem?

Propagation Formula

Indoor propagation loss formula (11n) *,

For d<16.5*ft*

 $Lp = -38 + 20 \log F + 20 \log d + Wall/Floor loss$ For d>16.5ft (Free Space formula)

F in MHz, d in feet

 $Lp = -38 + 20 \log F + 20 \log 16.5 + 35 \log (d/16.5) + Wall/Floor Loss$ Std. Dev 3-4dB (Shadow Loss)

*Erceg et al (2004) as per 11n, Channel Model B – Residential

- 10dB Outer Wall loss has been used in calculations
- No internal wall or floor losses used in calculations





Neighbors 3 houses down, and opposite houses within 150 feet have potential to overlap

Note: No internal wall losses, external wall loss only.

Submission



Neighbors 1 house down, and opposite houses within 150 feet have potential to overlap

doc.: IEEE 802.11-09/0762-00-00aa

Detached Houses

Woking, England

150 ft



12 Potential APs in range

Town Houses - Dense

Bleiswijk, The Netherlands



25 Potential APs in range

Submission

doc.: IEEE 802.11-09/0762-00-00aa

Terraced Houses

Leigh Park, Havant, England





16 Potential APs in range

Apartment Block Single Layout









Total within range = 28

NOTE: If each AP restricted operation to 54Mbps (11a/g), then overlaps reduce to 18 (or less)





Total within range = 28 + 25 = 53

NOTE: If each AP restricted operation to 54Mbps (11a/g), then overlaps reduce to 29 (or less)

Summary

- Examples used show maximum potential number of APs within range
 - Detached Houses 12
 - Terraced Houses 16
 - Townhouses 25
 - Single Layout Apartments 28
 - Double Layer Apartments 53

• Number of Channels

 - 2.4GHz
 20MHz
 3

 - 5GHz
 20 MHz
 24 USA, 19 Europe

 40MHz
 11 USA, 9 Europe

doc.: IEEE 802.11-09/0762-00-00aa

Channel Selection

Aside - Minimum number of Channels Apartments single layout

In fact, if one applied standard channel re-use to the Apartment single layout, 28 overlapping APs, <u>only 8 channels are actually required.</u>



Channel Selection Analysis Program

• A program was written in order to analyze what happens when each AP uses a Channel Selection scheme.

Program outline

- Randomly select an Apartment/House
- Scan the surrounding apartments/houses in range
 - First select channel(s) with least other APs
 - If more than one, then select channel with least total overlaps
 - E.g. Channel 1: 2 other APs, 1 overlap (one AP already sharing)
 - Channel 2: 2 other APs, 0 overlaps.
 - Selection will pick Channel 2
- Update each apartment/house with the number of other APs with which it is sharing
- The objectives are:
 - Determine how many channels are required to 'guarantee' zero or one overlaps
 - Investigate the overlap situation and "AP chains"
 - Use results to determine requirements for the OBSS solution

Detached Houses – 12 overlaps

			•												
		% Houses Assigned Channels													
Channels	10	20	30	40	50	60	70	80	90	100					
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000					
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000					
19	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000					
17	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000					
11	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000					
9	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000					

Channel Selection finds a clear channel – same result if Channel Only or Channel plus Overlaps selection

3 Channels does not work

Possibility Zero overlap

Percentage of Houses to be assigned:100 Probability of no overlaps: 0.1415 Probability of zero or one overlap: 0.5253 Probability of two overlap: 0.3551 Probability of three plus overlaps: 0.1195

Terraced Houses – 16 overlaps

Possibility Zero overlap

		% Houses Assigned Channels												
Channels	10	20	30	40	50	60	70	80	90	100				
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
19	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
17	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
11	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9998	0.9986				
9	1.0000	1.0000	1.0000	1.0000	1.0000	0.9997	0.9974	0.9880	0.9653	0.9144				

Possibility Zero or 1 overlap

		% Houses Assigned Channels												
Channels	10	20	30	40	50	60	70	80	90	100				
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
19	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
17	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
11	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				
9	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000				

doc.: IEEE 802.11-09/0762-00-00aa

Town Houses – 24 overlaps

	Possibilit	y Zero ove	erlap								
				% Apa	rtments As	signed Ch	nannels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	
19	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	
17	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	
11	1.0000	1.0000	1.0000	1.0000	1.0000	0.9997	0.9986	0.9850	0.9624	0.9146	
9	1.0000	1.0000	1.0000	1.0000	0.9988	0.9823	0.9437	0.8731	0.7917	0.6981	
	Possibilit	y Zero or	1 overlap								With 11 Channels
				% He	ouses Assi	gned Cha	nnels				with 11 Channels
Channels	10	20	30	40	50	60	70	80	90	100	100% chance that zero
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	Or single overlap
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	
19	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	
17	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	
11	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	K .
9	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9999	0.9992	0.9967	
	Possibilit	y 2 overla	р								
				% He	ouses Assi	gned Cha	nnels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
22	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
19	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
17	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
11	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.3% chance of a
9	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0001	0.0008	0.0033	Hiddon AP situation
	Possibilit	y 3+ overl	aps (AP Cl	hain)							
				% He	ouses Assi	gned Cha	nnels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
22	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
19	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
17	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
11	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
9	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	

9 0.0000 0.0000 Submission

Single Apartment Block – 28 overlaps

Possibility Zero overlap % Apartments Assigned Channels 30 70 80 100 Channels 10 20 40 50 60 90 1.0000 24 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 22 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 19 1.0000 17 11 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 0.9709 0.9161 0.8400 0.7574 1.0000 1.0000 1.0000 0.9985 0.9796 0.8227 0.7163 0.5912 0.4723 9 0.9227 Possibility Zero or 1 overlap % Apartments Assigned Channels With 17 Channels 100% chance that zero 10 20 30 40 50 60 70 80 90 100 Channels 1.0000 1.0000 1.0000 Or single overlap 24 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 22 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 19 1.0000 17 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 11 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 1.0000 0.9999 0.9998 0.9980 1.0000 1.0000 1.0000 0.9990 0.9948 0.9833 0.9579 9 1.0000 1.0000 1.0000 Possibility 2 overlap % Apartments Assigned Channels 10 20 30 40 50 60 70 80 90 Channels 100 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 24 0.0000 0.0000 0.0000 0.0000 22 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 19 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0 2% chance of hidden AP 17 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 With 11 channels 11 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0001 0.0002 0.0020 9 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0010 0.0053 0.1656 0.0386 Possibility 3+ overlaps (AP Chain) % Apartments Assigned Channels 20 30 70 80 90 Channels 10 40 50 60 100 0.0000 0.0000 0.0000 0.0000 24 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 22 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.02% chance of AP chain 19 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 17 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 With 9 channels 11 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 9 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000

Double Apartment Block – 53 overlaps

	Possibili	itv Zero o	verlap								L
	1	.,		% Apartn	nents As	sianed C	hannels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9952	0.9649	
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9999	0.9930	0.8497	0.8997	
19	1.0000	1.0000	1.0000	1.0000	1.0000	0.9987	0.9854	0.9308	0.9200	0.7402	
17	1.0000	1.0000	1.0000	1.0000	0.9990	0.9865	0.9239	0.8338	0.7164	0.6019	
11	1.0000	1.0000	0.9887	0.9198	0.7556	0.5786	0.4101	0.2609	0.1603	0.1017	
9	1.0000	0.9960	0.9233	0.7364	0.5067	0.3130	0.1801	0.0996	0.0600	0.0350	
	Possibili	ity Zero o	r 1 overla	ар							
				% Apartn	nents As	signed C	hannels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	99.27% chance of zero or one overlap
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	With 17 channels
19	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9997	0.9992	
17	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9999	0.9986	0.9927	
11	1.0000	1.0000	1.0000	1.0000	0.9968	0.9856	0.9460	0.8635	0.7342	0.5926	
9	1.0000	1.0000	0.9997	0.9964	0.9654	0.8748	0.7484	0.5682	0.4134	0.2850	
	Possibili	ity 2 over	lap								
				% Apartn	nents As	signed C	hannels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.73% chance of 2 overlaps with 17 channels
22	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
19	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0003	0.0008	
17	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0001	0.0014	0.0073	
11	0.0000	0.0000	0.0000	0.0000	0.0032	0.0144	0.0536	0.1312	0.2411	0.3444	
9	0.0000	0.0000	0.0003	0.0036	0.0343	0.1208	0.2276	0.3535	0.4140	0.4157	
	Possibili	ity 3+ ove	erlaps	~ ~ ~							1
Channela	10		20	% Apartn	nents As	signed C	nanneis			100	With 17 Channels no cases of 3 overlaps
Channels	10	20	30	40	00	00	70	80	90	100	
24	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
40	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
19	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	r
11	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0004	0.0053	0.0247	0.0630	
9	0.0000	0.0000	0.0000	0.0000	0.0003	0.0044	0.0240	0.0783	0.1727	0.2993	
~	0.0000	0.0000	0.0000	0.0000	0.0000	0.0011	0.01	5	0.1121	0.2000	

Submission

Double Apartment Block – 53 overlaps –Not using overlap selection

				% Apartn	nents As	signed C	hannels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9959	0.9650	
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9999	0.9956	0.9625	0.8929	
19	1.0000	1.0000	1.0000	1.0000	1.0000	0.9992	0.9834	0.9247	0.8406	0.7364	
17	1.0000	1.0000	1.0000	1.0000	0.9998	0.9821	0.9305	0.8199	0.7038	0.5862	
11	1.0000	1.0000	0.9957	0.9116	0.7570	0.5763	0.4222	0.2953	0.2073	0.1390	
9	1.0000	0.9995	0.9273	0.7253	0.5255	0.3508	0.2249	0.1373	0.0866	0.0510	
	Possibil	ity Zero o	r 1 overla	ар							
				% Apartn	nents As	signed C	hannels				Com
Channels	10	20	30	40	50	60	70	80	90	100	
24	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9997	
22	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9999	0.9997	0.9976	Chan
19	1.0000	1.0000	1.0000	1.0000	1.0000	1.0000	0.9999	0.9990	0.9959	0.9915	ofov
17	1.0000	1.0000	1.0000	1.0000	1.0000	0.9999	0.9991	0.9956	0.9858	0.9687	
11	1.0000	1.0000	1.0000	0.9976	0.9862	0.9514	0.8943	0.7980	0.6847	0.5720	
9	1.0000	1.0000	0.9983	0.9810	0.9310	0.8304	0.6949	0.5523	0.4211	0.3200	
	Possibil	ity 2 over	lap								
				% Apartn	nents As	signed C	hannels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0003	
22	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0001	0.0003	0.0023	
19	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0001	0.0001	0.0041	0.0084	
17	0.0000	0.0000	0.0000	0.0000	0.0000	0.0001	0.0009	0.0044	0.0138	0.0307	
11	0.0000	0.0000	0.0000	0.0024	0.0132	0.0470	0.0997	0.1821	0.2654	0.3259	
9	0.0000	0.0000	0.0017	0.0186	0.0657	0.1516	0.2574	0.3307	0.3662	0.3582	
	Possibil	ity 3+ ove	erlaps								
		_		% Apartn	nents As	signed C	hannels				
Channels	10	20	30	40	50	60	70	80	90	100	
24	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
22	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
19	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0001	
17	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0004	0.0006	
11	0.0000	0.0000	0.0000	0.0000	0.0006	0.0016	0.0060	0.0200	0.0498	0.1020	
1	0.0000	0.0000	0.0000	10.0004	0.0022	0.0400	0.0477	0 1170	0.2427	0.2240	1

pare to previous slide

nel Selection using the number erlaps is better

Channel Selection Analysis (08/1470r4)



Submission

doc.: IEEE 802.11-09/0762-00-00aa

Number of Channels? What about 40/20 NHz?

40/20 MHz Example

Run the Channel Selection program

1. 40 MHz scenario

- 100 Apartments (10 floors of 10)
- 28 potential apartments within range
- 9 Channels, (40MHz channels)

2. 20 MHz Scenario

- 100 Apartments (10 floors of 10)
- 28 potential apartments within range
- 17 Channels, (20MHz channels)

In each case look at the channel sharing results

May 2009



100 Apartments (10 floors of 10)28 potential apartments within range9 Channels, (40MHz channels)

Result:

- 58 Apartments not sharing
- 34 Apartments share with one other
- **3** Apartments share with **2** others
- 5 Apartments share with 4 others

Double apartment block is worse with 54 potential overlaps

Submission

20/40MHz channels

Assume that a sharing 40MHz channel, dropped back to 20MHz

- The 1 : 1 share can be two independent 20MHz channels
 - e.g. L : U (Lower 20MHz, Upper 20MHz)
- The 1 : 2 : 1 chain can be independent 20MHz channels - e.g. 2L : 2U : 2L
- The 1 : 2 : 1 : 2 : 1 can be also be independent 20 MHz channels – e.g. 7L : 7U : 7L : 7U : 7L

Now:

- 58 Apartments using 40MHz channel, not sharing
- **42** Apartments using 20MHz channels, not sharing Compare with (40MHz)
 - 34 Apartments share with one other
 - 3 Apartments share with 2 others
 - 5 Apartments share with 4 others

<u>May 2009</u>

OVERL/	APS			1					
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0
CHANN	EL #		-						-
6	3	4	1	2	6	3	7	1	4
1	5	9	7	12	11	5	8	10	2
4	2	8	3	4	1	2	6	3	7
7	11	6	5	10	9	12	11	5	8
3	12	4	2	8	7	3	13	4	9
6	1	7	11	6	5	10	1	2	11
4	5	8	3	12	4	8	7	3	6
9	2	10	1	9	2	6	11	10	1
3	11	6	4	7	3	12	4	9	5
1	7	9	2	8	5	0	1	8	2

100 Apartments17 Channels, 20MHz channels

Result: 100 Apartments not sharing

40MHz Sharing compared to 20MHz

• **Theoretical** (Figs 8.1 and 8.5 "Next Generation Wireless LANS", Perahia and Stacey)

2 x 2

- 40MHz channel, max throughput 91Mbps and 210Mbps (RIFS + BAR)
- 20MHz channel, max throughput 62Mbps and 100Mbps (RIFS + BAR)

3 x 3

- 40MHz channel, max throughput 85Mbps and 300Mbps (RIFS + BAR)
- 20MHz channel, max throughput 69Mbps and 150Mbps (RIFS + BAR)

Assuming perfect sharing, 20 MHz throughput is better or same as shared 40MHz throughput

- Practical (two overlapping networks, AP and 1 STA)
 - Individually (out of box)
 - 40MHz Channel BSS A: 77-80Mbps BSS B: 58Mbps
 - 20MHz Channel BSS A: 40-50Mbps BSS B: 60Mbps
 - Sharing
 - 40MHz Channel BSS A: 77-80Mbps BSS B: 2-7Mbps

DCF Sharing is not ideal, it seems!

Observation

If sharing with an overlap of 2, then definitely in everyone's interest to drop back to 20MHz channel

If sharing with an overlap of 1, then could consider sharing IF you consider that sharing 40MHz is better than an independent 20MHz

BUT

If you believe that in practice, devices will share equally on 40MHz, I have a bridge I want to sell you.

Conclusion

Equal sharing using DCF and EDCA is not always true in practice. Smaller number of channels results in OBSS "Chains"

Drop back to 20MHz channels when no clear 40MHz channel is available.

Extend this to higher bandwidths?

doc.: IEEE 802.11-09/0762-00-00aa



PROPOSED "QLOAD" ELEMENT

QLOAD ELEMENT



Overlap

- QAP indicates the number of other QAPs with which it is sharing and indicates the size of the OBSS graph:
 - Zero indicates QAP has no other QAPs on the same channel within range
 - 1 indicates already sharing with one other QAP
 - 2 indicates already sharing with two other QAPs
 - etc

The QAP is advertising the overlap to other QAPs who may be considering sharing. This parameter should be included in the Channel Selection procedure in order to select the best channel (08/1470r4)

Distance

- Distance is set to 0 for Self
- If QAP ID Directly visible to the QAP Self, then "Distance" is set to 1
- If not directly visible to the QAP Self, then "Distance" is set to 1 plus the value reported for that QAP ID in the QAP that is directly visible
- Any QAP with Distance" > 2 is not recorded in QLoad Element

QAP ID

- First octet = random number (0 to 255)
- Second octet = octet 6 of MAC Address
- Once established, QAP ID is not changed
- Enables a QAP to indentify its QLoad in other QLoad elements

QLoad Self

There are three methods for the QAP to build QLoad Self:

- 1. QSTAs in the BSS may send a TSPEC (ADDTS) with Inactivity Interval set to 0 (or 1) for instant timeout
 - By sending in a TSPEC the STA has the QAP commit, in advance, medium time for the STA

2. QAP notes and adjusts for new TSPECs from QSTAs

- If accepted, "QLoad Self", and also "QLoad Total" are adjusted only when the QSTA submits the ADDTS
- Chance that ADDTS is denied as QSTA did not reserve medium time in advance

3. In response to TSPEC Requirements Request

- QAP request STAs to indicate or confirm their TSPECs
- Used by QAP to 'clear house' or initially set up Q Load.

The QAP is advertising its own potential QoS load to other QAPs who may be considering sharing

TSPEC Requirement Request Response

- Request from QAP to a particular STA
- Two types of Request:
- 1. Send All TSPECs (ID 0)
 - Effectively all previous (if any) TSPECs are deleted, need to set them up again
- 2. Confirm TSPECs (ID 1)
 - Confirm which TSPECs are still required
 - TSID plus Direction defines TSPEC



QLoad MEAN and STDEV

For detailed explanation see 09/0496r2

MEAN and STDEV is estimated from the individual TSPECs:

MEAN	$\mu = \Sigma MEANi$
STDEV	$\sigma = 0.25 \text{ sqrt} \{ \Sigma (MAXi - MINi)^2 \}$

MEAN	$\mu_{tot} = \Sigma MEANi$
STDEV	$\sigma_{tot} = sqrt(\Sigma \sigma_i^2)$

Total Traffic Requirement can be estimated:

- 1. MAX traffic = $\mu_{tot} + 2 \sigma_{tot}$
- 2. 90% Traffic = μ_{tot} + 1.3 σ_{tot}
- 3. 80% Traffic = $\mu_{tot} + 0.83\sigma_{tot}$

EDCA Overhead – Capacity drops with # streams

As number of video streams increases, the contention also increases. In order to keep latency low the capacity of the Channel is decreased. Maximum throughput on (shared) channel decreases as number of video streams increases



NOTE: Above graph is simulation for independent streams. Downlink streams from QAP may be better due to queuing at the AP **Total Allocation MUST take account of the number of streams** *Note: This is also for Admission Control on each QAP*

QAP Priority Streams

- Number of EDCA Priority Streams, AC_VO and AC_VI
- Used to estimate "EDCA Bandwidth Factor"
- EDCA Bandwidth Factor = 1 + 0.05 N (approx; keep it simple, see 09/0497)
 - Where N = Number of streams
 - Example:

4 streams Effective Bandwidth Factor = 1.2

Four 5.5Mbps streams will require 1.2 x 4 x 5.5 = 26.5Mbps

Access Fraction and Access Factor

Access Fraction

- Total actual admitted time and/or scheduled time expressed as a fraction of 32us/sec rounded down to 1/256
- Access Factor
 - Total Traffic Requirement of self plus all other visible QAPs.
 Expressed as a fraction that may be greater than 1
 - Calculated as follows:
 - Sum the individual QLoads of all QAPs in the QLoad element as a composite stream
 - Calculate the EDCA Bandwidth Factor from the total number of Priority Streams in the visible QAPs (Distance 0 and 1)
 - Multiply the two to obtain the "Access Fraction"

Sharing

- If the Access Factor is >1, then there is a potential overallocation
 - Hopefully QAPs should avoid this in the Channel selection process

• Sharing Scheme

- QAPs should examine their QLoad Element in order to determine the maximum "Access Factor" being reported. This maximum value is then used to determine the allocation limit for that QAP in order not to cause over-allocation in other QAPs that are overlapping,
- Using the Access Fractions (actual "live" traffic), Access Factor and QLoad self, a decision can be made whether to admit a new request.
- Rules could be recommended in informative text.

HCCA Peak

- The total HCCA TXOP requirement for the QAP, expressed in 32us/sec.
 - The sum of all the HCCA Peak values is the "HCCA Access Factor"
 - If HCCA Access Factor > 1sec then potential for TXOP overallocation
 - HCCA TXOPs can sum to "1" independent of EDCA Medium Time allocations, as TXOPs terminate immediately when no more data

Medium Time Allocations - Sharing

- It is important to understand how the AP allocates the actual Medium Times in responses to TSPECs and checks that it has not exceeded its 'limit'
- 1. In response to each TSPEC the AP <u>allocates the Medium Time or</u> <u>TXOP (HCCA) that corresponds to the peak traffic</u>
- 2. When allocating an additional Medium Time or TXOP, the AP must calculate what the <u>composite</u> stream would be, and check that this <u>composite</u> medium time does not exceed the limit
- 3. It is this <u>composite time</u>, that is advertised in the <u>Access Fraction</u>
 - The actual sum of the Medium Times and TXOPs will be greater than the composite time, but EDCA only uses what it needs, and hence the statistical nature of the streams causes the composite time to be the maximum of what is actually being used. Similarly HCCA TXOPs terminate when no more data.
 - Allocated HCCA TXOPs cannot exceed "1"

HCCA

• Two schemes are under consideration for HCCA TXOPs scheduling

1. HCCAOP Advertisement element

- This is based upon the MCCA Element used in 11s. Each QAP advertises its TXOP and the others then avoid it
- Suffers from clock drift and 'gaps' in the schedules
- Expandable to any size OBSS and hidden QAPs

2. Supervisor QAP

- Supervisor hands off the schedule to each QAP using Wireless DS to communicate AP to AP
- Efficient
- Limited to visible QAPs

doc.: IEEE 802.11-09/0762-00-00aa



CONCLUSIONS

• OBSS is only a real problem for QoS

- If 802.11 wants to provide guaranteed performance for QoS applications, then something has to be done
- OBSS is a real possibility, even at 5GHz.
- The number of available channels is very important
 - Dropping back to lower channel bandwidth, and increasing the channel pool has distinct advantages
- Channel selection using inputs such as "Overlap" and "QLoad" is very desirable
- The proposed QLoad element provides all the necessary information to enable 'good' sharing for QoS networks.