

EtherCAT Relay function

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Disclaimer

This presentation should be considered as the personal views of the presenters and not as a formal position, explanation, or interpretation of IEEE and ETG.

More details can be found in the presentation from 2017:

<https://www.ieee802.org/1/files/public/docs2017/liaison-ETG-streamAdaption-1117.pdf>

The website of the EtherCAT Technology Group (ETG) provides further information about EtherCAT: https://ethercat.org/en/tech_group.html

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IEEE 802 challenge for high speed applications

- Efficiency: low byte count (8 bytes) needed vs. 84 octets minimum for Ethernet
- Delay: cable delay of fieldbusses vs. store and forward/bridging (passive media used for first solutions)
+ interfering traffic

➔ **Overall efficiency 3%**
(with 128 octets interfering traffic)

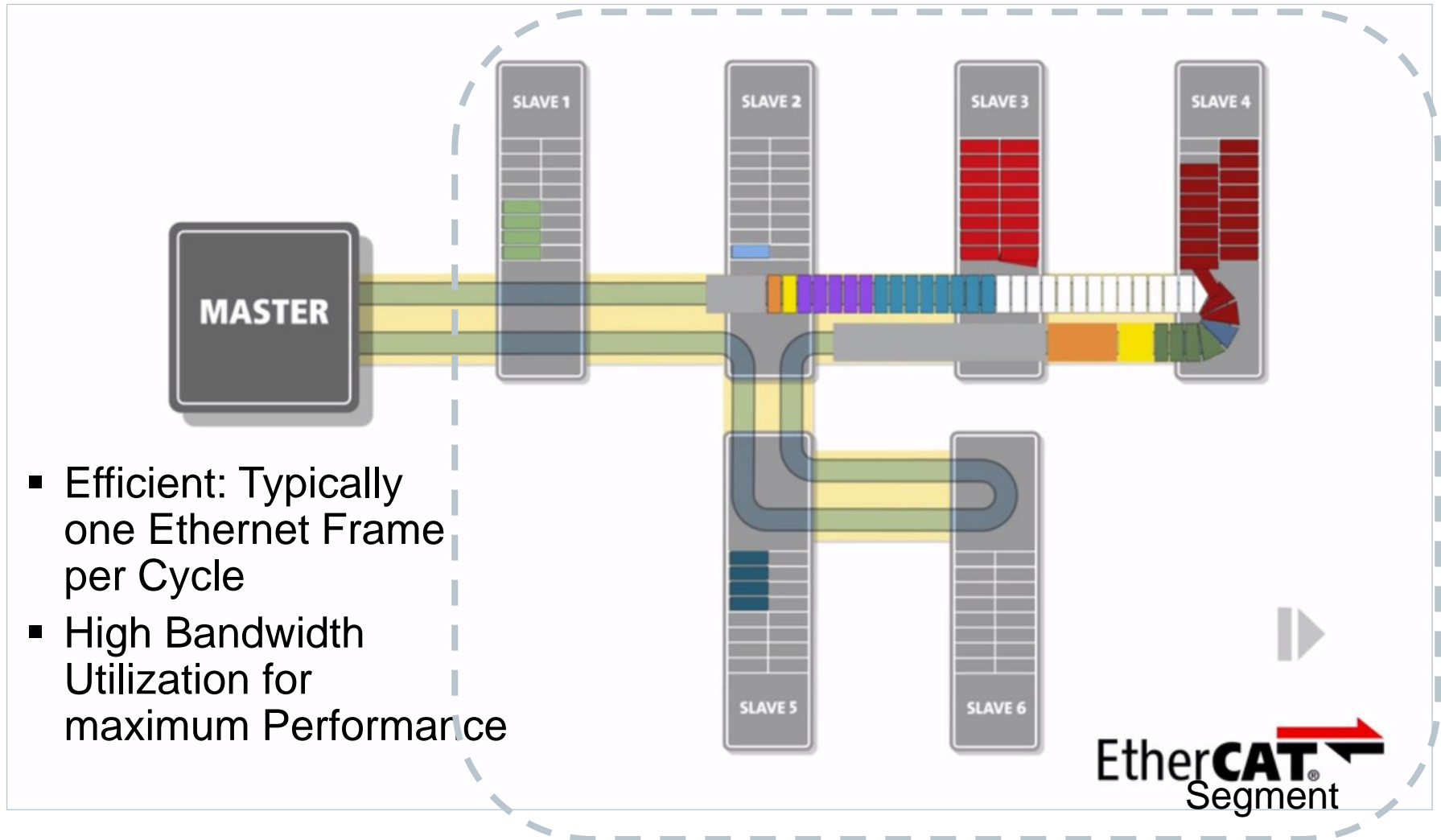
- This leads to the **EtherCAT** approach

EtherCAT®



- Efficiency **Shared frame** instead of individual frame
➔ performance improvement: overhead 50 Bytes instead of 750/1500 ... in a network of 10/20 I/O stations
- **Processing on the fly** with topological relay function (automatic)
Instead of address based forwarding
➔ performance improvement: 0,6 μ s instead of >3 μ s (7 μ s/store&forward)

Functional Principle: Ethernet “on the fly”



*Master/Slave shall be replaced by MDevice/SubDevice

Animation available as EtherCAT Functional Principle (2D) on <https://www.youtube.com/watch?v=z2OagcHG-UU>

Preconditions for the segment forwarding

- Same Data Rate within the segment (typical 100 Mb/s – 1000 Mb/s possible)
 - The MDevice will send out frames and receive frames
 - Each SubDevice in the Segment just relays frames to the next port
 - The forwarding acts as a unidirectional relay from port to port
 - Receive Port 0 connected to a virtual Port to allow a DLL entity to put information in the frame and get information „on the fly“
 - The forwarding takes place on any open port
 - The forwarding already starts with the Preamble
 - If there is only one port open the relay function will send it back at the same port
 - Each device that is connected will receive the frame
 - The flow of data in a segment forms a logical ring of all devices connected
- ➔ The ports of a node are ordered as a ring with Port0 as first entry
- ➔ If the main structure is a line, the exits of the line shall be in between Port0 and Port(last) for easier handling and diagnostic (no technical reason)

Data Flow - Example

Backbone 1

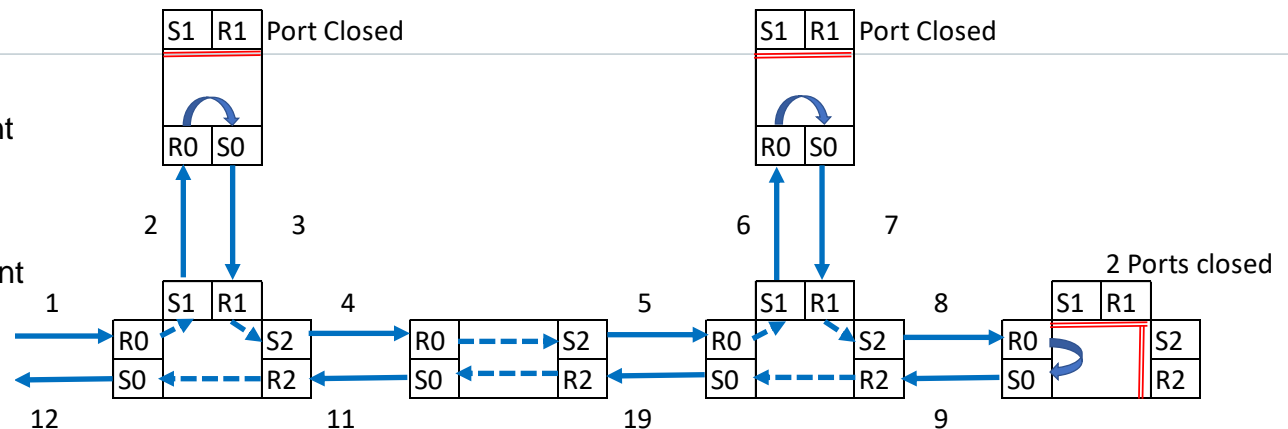
1st Sub-segment

Backbone 2

Backbone 3

2nd Sub-segment

Backbone 4



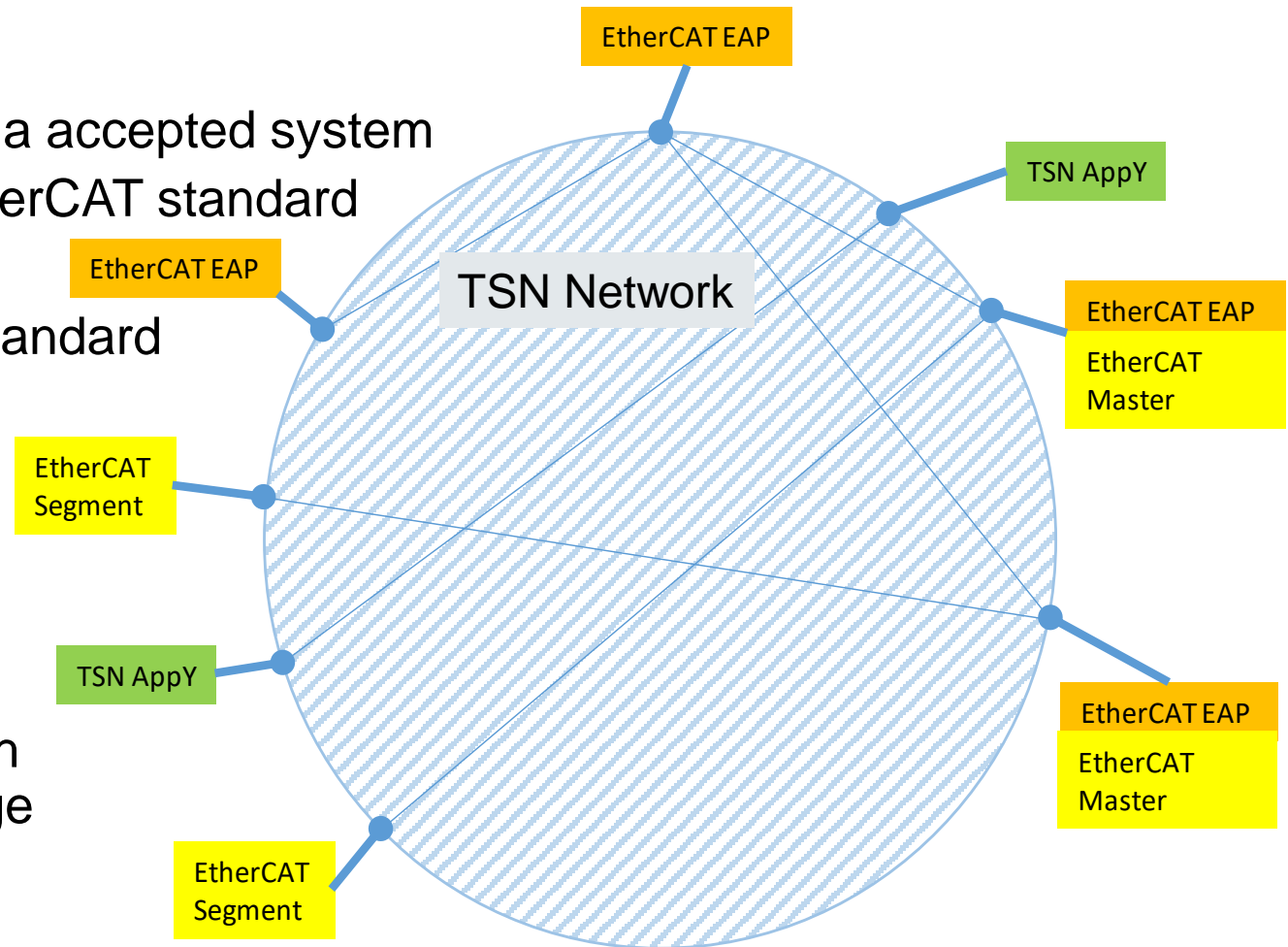
Use of Open mode → TSN

- Idea

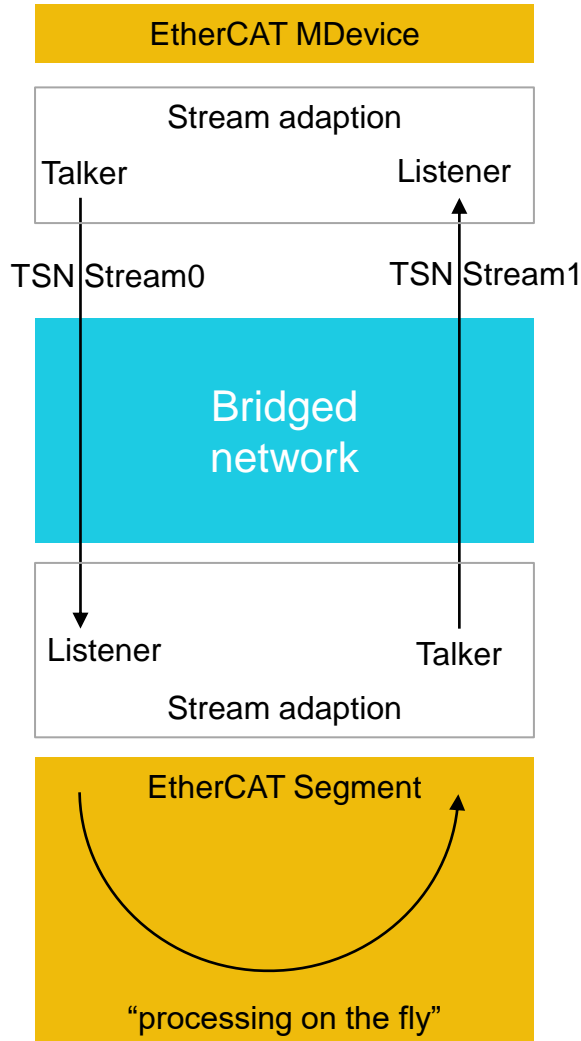
- Never touch a accepted system
- Maintain EtherCAT standard
- Use 802.1 standard

- Compatibility with other 802.1 usage

- Specifies adaptation

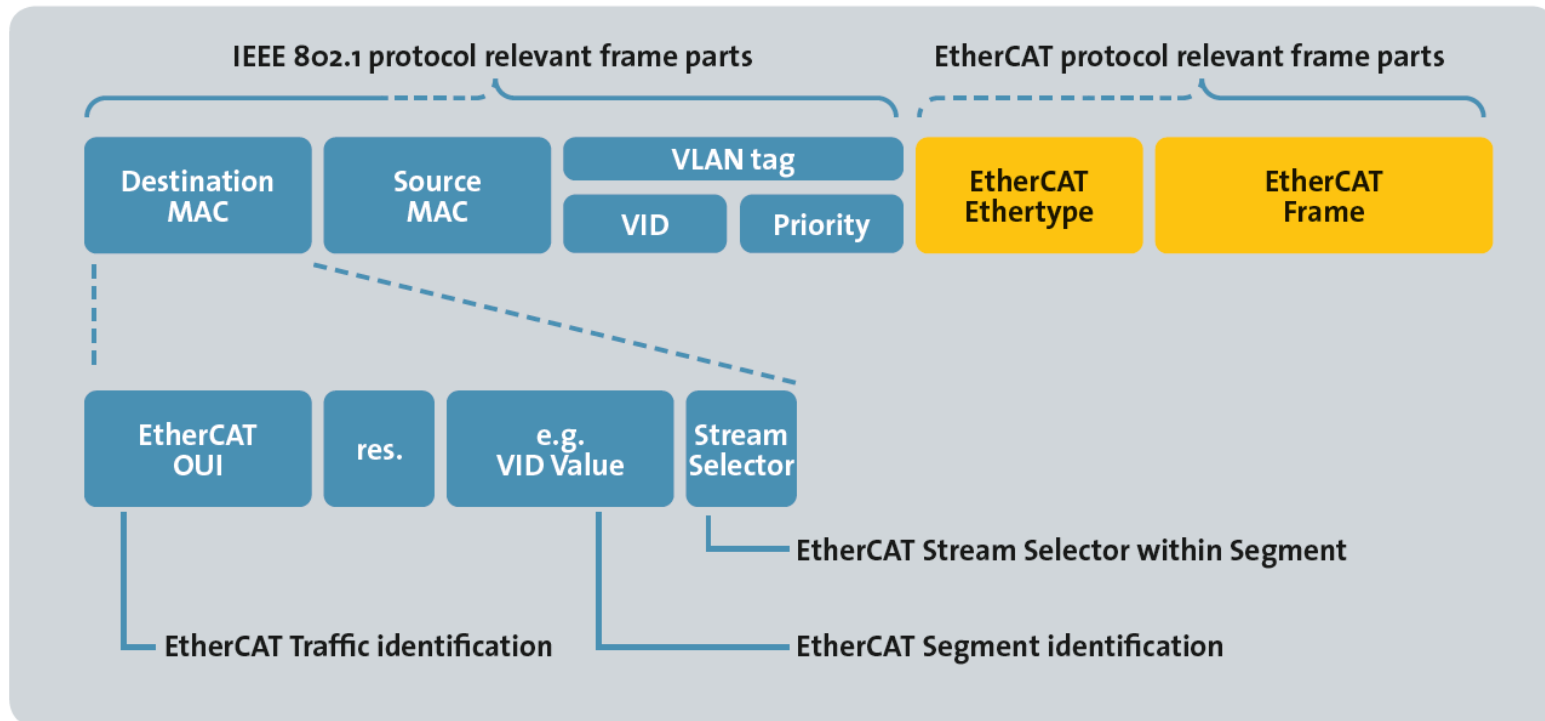


Stream Adaptation: Details



- Always a pair of streams is set up
- Minimum one pair, but more might be set up, e.g.
 - One for cyclic
 - One for acyclic (strict priority)
 - for additional transfers
- Traffic class for pair of stream always the same
- Maintain Traffic Class (VLAN Prio)
- Maintain length (EtherCAT Rx/TX frame length identical)
- **Unique Stream Identification required(!)**

Protocols use different fields



Open Mode

- EtherCAT segment corresponds to an Identifier (“VID”)
- MAC addresses for stream identification (StreamDA, StreamSA) constructed of
 - OUI, (V)ID, Stream selector
- Multicast DAs are possible as stream MAC for TSN Networks
Null Stream Identification combined with Source MAC and VLAN Stream Identification according to Table 6-1 of 802.1CB
- Unicast DAs are possible for streams to the segment
Address change for backward direction to avoid multicast scans in the MDevice and enables address learning in network